

Claims

What is claimed is:

1. A method of priority queue dispatching in a data processing system, wherein:
a critical section of code for a task is when the task has at least one of a plurality of locks locked; and
said method comprises:

- A) dispatching a first task for execution at a temporary dispatching priority higher than a standard dispatching priority for that first task whenever the first task is in a critical section of code for that first task, wherein the critical section of code for that first task is a result of that first task having locked a first lock without having unlocked the first lock; and
- B) dispatching the first task for execution at the standard dispatching priority for that first task whenever the first task is not in the critical section of code for that first task.

2. The method in claim 1 wherein:
the method further comprises:

- C) queuing a second task in a first FIFO queue when the second task attempts to lock the first lock and fails, wherein:
the FIFO queue comprises a set of tasks waiting to lock the first lock; and
step (C) comprises:
 - 1) placing the second task at an end of the first FIFO queue;
 - 2) comparing the temporary dispatching priority for a third task that is ahead of the second task in the FIFO queue to the temporary dispatching priority of the second task; and
 - 3) setting the temporary dispatching priority of the third task to the temporary dispatching priority of the second task when the temporary dispatching priority of the second task is determined in substep (2) to be greater than the temporary dispatching priority of the third task.

- 1 5. The method in claim 4 wherein step (C) further comprises:
2 4) testing whether the first task is in a second FIFO queue awaiting a
3 chance to lock a second lock; and
4 5) upgrading the temporary dispatching priority of a third task that
5 has a second lock locked when the first task is determined in
6 substep (4) to be in the second FIFO queue awaiting a chance to
7 lock the second lock, wherein substep (5) comprises:
8 a) comparing the temporary dispatching priority for the
9 third task to the temporary dispatching priority of the
10 second task; and
11 b) setting the temporary dispatching priority of the third
12 task to the temporary dispatching priority of the second
13 task when the temporary dispatching priority of the
14 second task is determined to be greater than the
15 temporary dispatching priority of the third task.
- 1 6. The method in claim 5 wherein:
2 the method further comprises:
3 D) comparing a highest temporary dispatching priority of any task in
4 the first FIFO queue to the temporary dispatching priority of the
5 first task after the first task unlocks the second lock; and
6 E) setting the temporary dispatching priority of the first task to the
7 highest temporary dispatching priority of any task in the first
8 FIFO queue when the temporary dispatching priority of the first
9 task is determined in step (D) to be greater than the highest
10 temporary dispatching priority of any task in the first FIFO
11 queue.

8. The method in claim 1 wherein:
the temporary dispatching priority for the first task is determined
dynamically and adjusted dynamically depending on the
temporary dispatching priority for other tasks behind the first
task in a first FIFO queue associated with the first lock.

1 9. The method in claim 1 wherein:
2 the temporary dispatching priority for the first task is a system wide
3 value.

1 10. The method in claim 1 wherein:
2 the temporary dispatching priority for the first task is a value
3 associated with the first lock when the first lock is a last lock
4 that the first task has attempted to lock.

- 21 -

